

**6331 - Algorithms, Autumn 2016, CSE, OSU**

**Homework 3**

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**Problem 1.**

- (a) Starting with an empty binary search tree, give a sequence of  $n$  insertion operations that result in a tree of height  $\Omega(n)$ .
- (b) Starting with an empty binary search tree, give a sequence of  $n$  insertion operations that result in a tree of height  $O(\log n)$ .

**Problem 2.** In Section 12.4 of the book it is shown that the expected height of a randomly built binary search tree on  $n$  distinct keys is  $O(\log n)$ . That is, if we start with an empty tree, and we insert  $n$  distinct elements in a *random* order, then the expected height of the tree is  $O(\log n)$ . Here, by “random order” we mean an order chosen uniformly at random from the set of all possible orderings of the  $n$  keys. Also note that since the ordering is random, the height of the resulting tree is a random variable. In other words, the above result states that the expectation of this random variable is  $O(\log n)$ .

Use the above fact to construct a randomized algorithm for sorting  $n$  elements, with expected running time  $O(n \cdot \log n)$ .

Hint: You don’t need to use the proof of the above statement about the expected height from Section 12.4. It is enough to assume that the statement is true.